

Context-Free Grammar - Exercise

Problem:

Write a context-free grammar for $L = \{a^n b^m c^k : n+2m=k, n \geq 0, m \geq 0\}$.

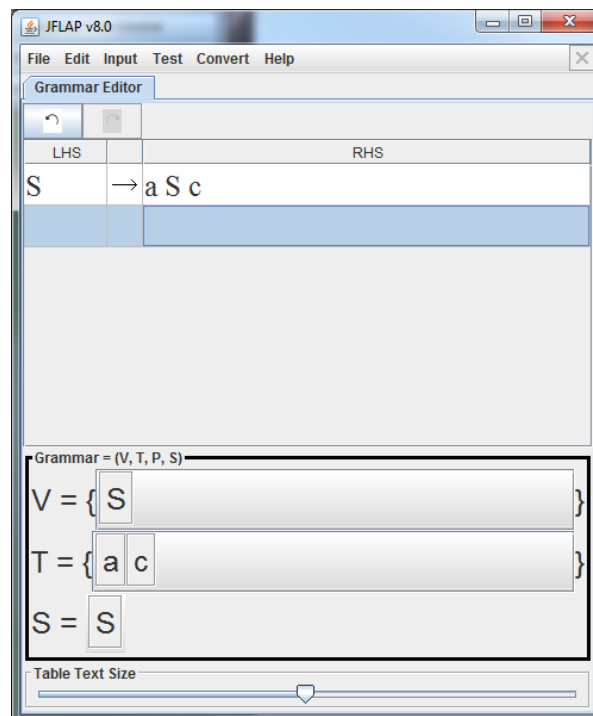
Solution:

First, we examine the kinds of words produced by this set. One way to do that is to tabulate the different values of n , m , and k and list the corresponding words for the different values.

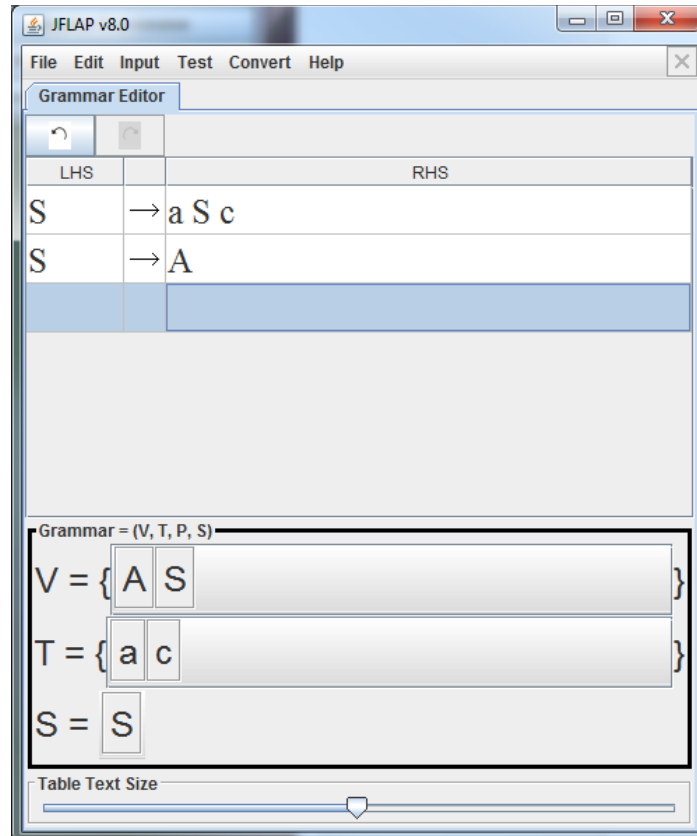
n	m	k	w
0	0	0	λ
0	1	2	bcc
1	0	1	ac
1	1	3	abccc
1	2	5	abbccccc
			...

We see that for each a in the word, there is a c produced. For each b in the word, there are two b 's. Also, a 's, b 's, and c 's appear in that sequence; i.e., $ccbba$ is not accepted.

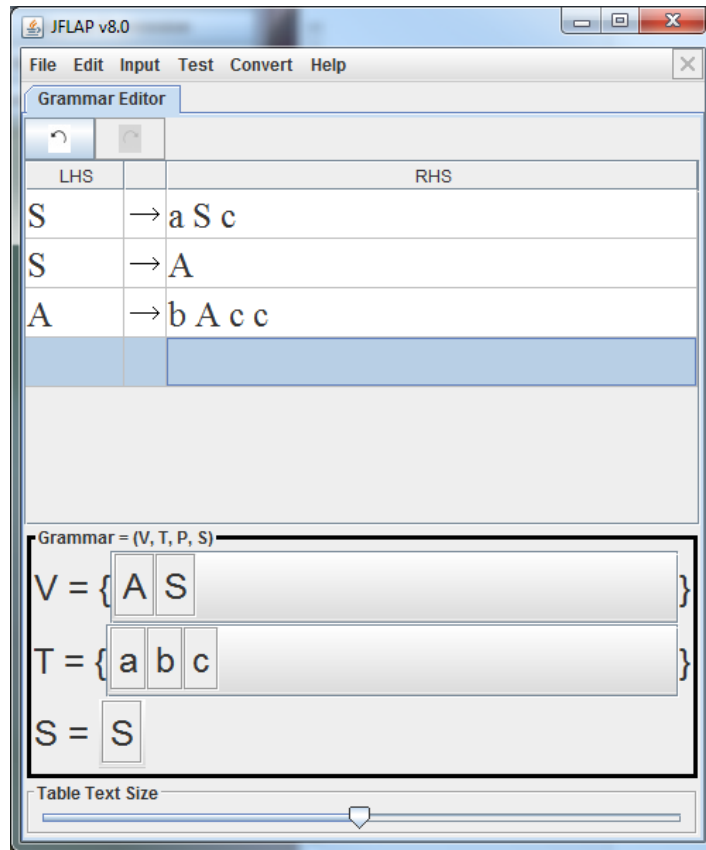
Open JFLAP and click on *Grammar*. Notice that the resulting window has a LHS column and a RHS column. Each row refers to one rule in the context-free grammar to be created. Using S as the start symbol, fill the first row with $S \rightarrow aSc$ like so:



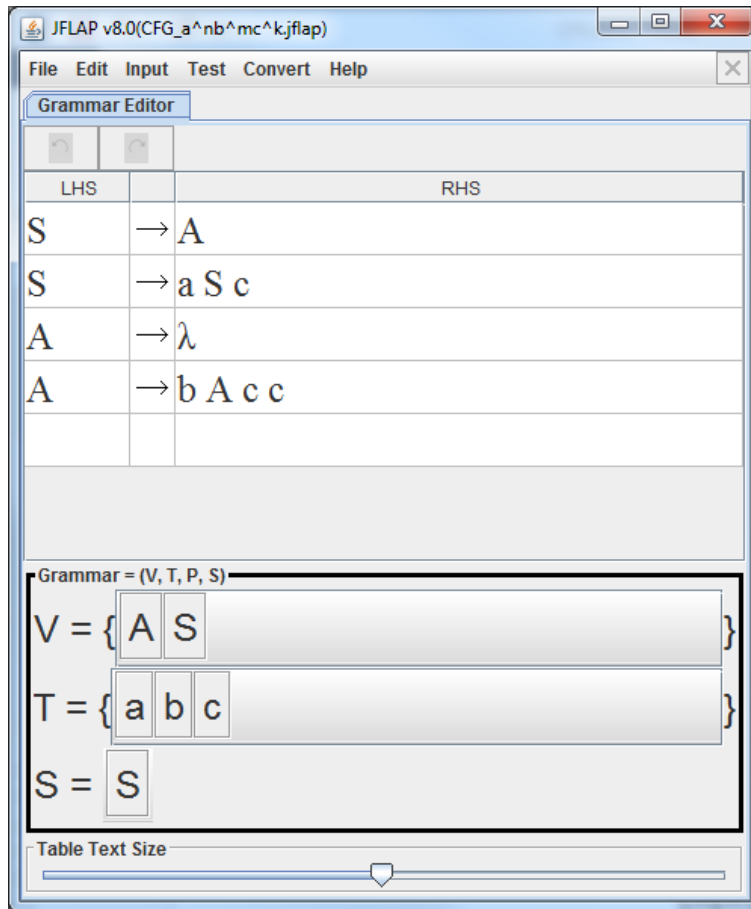
The S in the RHS allows the rule to become recursive. This rule states that for every a created, there is a c at the end of the word created. Next, we introduce a production rule to allow b's to be created. Note that once b's are being produced, we cannot create any more a's. We do this by adding a second rule, $S \rightarrow A$.



Next, we work on letting b's be created making sure that there are two c's for every b, $A \rightarrow bAcc$. This rule is also recursive allowing an infinite number of b's to be created.



Last, we add a rule to terminate the derivation, $A \rightarrow \lambda$.



That completes the grammar. Next, we test for grammar type with *Test > Test for Grammar Type*. JFLAP reports that it is a context-free grammar.

Lastly, we may test an input string. Click *Input > Brute Force Parse* and enter *abbccccc* as the input. Click *Complete*. JFLAP reports this string as *accepted*. The brute parse table also shows the 5 derivation levels and the progression from S to the test input.

JFLAP v8.0

File Help

Grammar Editor Brute Force Parser

Input:

Input accepted! Change view to see derivation!

S	→ A	Brute Parse Table	
S	→ a S c	Level	Total Nodes
A	→ λ	1	2
A	→ b A c c	2	6
		3	8
		4	10
		5	11
			Current Derivations
			[A, a S c]
			[λ, a A c]
			[a b A c c c]
			[a b b A c c c c c]
			[a b b c c c c c]

Grammar = (V, T, P, S)

V = { }

T = { }

S =

Table Text Size